

## A Personalized Hierarchical Feature Of Cyber Text Based Proxy Encryption Using Mobile Computing

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### ABSTRACT

*The Proxy encryption strategies with respect to comfy cloud records and its utility. There are various encryption schemes that offer protection and get right of entry to manage over the community. Proxy encryption enables the semi-depended on proxy server to encrypt .The encryption is accomplished with-out the server being capable of decrypt. The function primarily based definitely proxy encryption (ABPRE) scheme is one of the proxy encryption, which. ABPRE extending the conventional proxy encrypted and attributes an critical characteristic.*

*In ABPRE Proxy encryption strategies with recognize to cozy cloud data and its utility. there are various encryption schemes that offer protection and get admission to manipulate over the network. Proxy encryption permits the semi-relied on proxy server to re-encrypt the legal consumer simply makes use of his very personal mystery key to decrypt the encrypted statistics, and he doesn't need to shop an extra decryption key for decoding; The touchy records cannot be observed to the proxy in encryption, and the proxy most effective complies to the facts owner's command. We achieve this motive by means of first combining the Hierarchical based totally Encryption (HIBE) system.*

*Keyword: Attribute-based proxy re-encryption, cloud computing, data sharing.*

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### I. INTRODUCTION

Cloud computing is an rising financial and computing paradigm with the development of internet generation. Diverse application offerings can be furnished to fulfill users' requirements through the cloud computing . considered one of cloud computing software services generally used is data storage and users can outsource their storage and shop their touchy statistics inside the cloud. Migrating facts from the consumer side to the cloud gives tremendous comfort to customers, seeing that they can get admission to records within the cloud anytime and everywhere, using any device, without being concerned about the capital funding to set up the hardware infrastructures.

The data proprietor outsources a hard and fast of records to the cloud. every piece of statistics is encrypted before outsourcing. The records owner is chargeable for figuring out the get entry to shape for each statistics, and dispensing user characteristic mystery keys (UAKs) corresponding to person attributes to each person. Attribute based totally Encryption (ABE) presents regular encryption and extra access manage feature. ABE is more green, bendy and suitable than other cryptographic strategies and may be a light-weight security solution for net offerings . The cloud service provider directs a cloud to provide facts storage provider. facts owners encrypt their data documents and keep them within the cloud for sharing with facts customers. To touch the shared records files, facts customers down load encrypted information documents in their hobby from the cloud after which decrypt them. every information proprietor/client is managed by means of a domain influence. a site authority is directed by its figure area authority or the believed authority. facts owners, domain authorities, records purchasers, and the conditioned authority are prearranged in a hierarchical way.

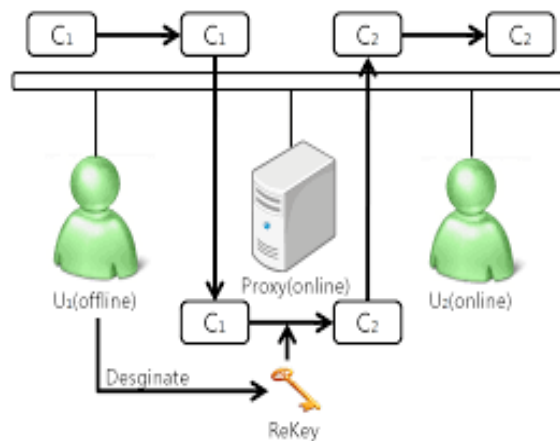
**II. LITERATURE SURVEY**

Cloud service carriers reveals out the access manage mechanisms for data at the cloud. access control is a way that restricts, denies, or lets in get entry to to device. in the cloud, facts protection is crucial to guard in opposition to inside assault, denial of provider attack, and collision assault[1]. An characteristic based totally encryption scheme turned into introduced by means of Sahai and Waters in 2005 and the purpose is to provide security and get admission to manage [2] . one in every of such technique is Proxy Re-encryption(PRE) approach[3]. In ABE method, the records is stored on the storage server in an encrypted form at the same time as extraordinary users are nonetheless allowed to decrypt different pieces of facts as consistent with safety coverage[3]. This successfully removes the need to rely on the storage server for preventing unauthorized statistics access.

The Hierarchical characteristic-based totally Encryption (HABE) is derived by way of Wang et al The HABE version includes a Root master (RM) that corresponds to the 1/3 depended on party (TTP),more than one domain Masters (DMs) in which the pinnacle-stage DMs correspond to multiple organization users, and severa users that correspond to all personnel in an enterprise[5][6]. This scheme used the property of hierarchical era of keys in HIBE scheme to generate keys. the principle purpose of cloud storage device is to at ease the information itself in this sort of manner that even in the occasion of a successful assault[7]. To offer confidentiality for messages in storage servers, a person can encrypt messages by way of a encryption technique to encode and shop messages [8][9].

**III. Attribute-based Proxy Re-encryption Algorithm (ABPRE)**

The based proxy re-encryption scheme. This encryption scheme guarantees facts confidentiality and high-quality gain get admission to manage .attribute-based totally proxy re-encryption (ABPRE) was brought by and enriched by way of with diverse capabilities. but, those solutions do no longer support the function of keyword search on encrypted records. the solution in this paper may be seemed as an extension to ABPRE with the feature of keyword seek on encrypted information.



**Fig1.1: SYSTEM MODEL**

**3.1 Proxy Re-Encryption (PRE)**

Allow us to illustrate the incentive of the PRE scheme by the subsequent example: Alice receives emails encrypted underneath her public key PKA thru a semi depended on mail server. while she leaves for vacation, she desires to delegate her e mail to Bob whose public secret's PKB, but does no longer need to proportion her mystery key SKA with him. The PRE scheme lets in Alice to offer a PRE key B->RKA to the mail server, with which the mail server can convert a encrypted that is encrypted underneath Alice's public key PKA into every other encrypted that may be decrypted via Bob's mystery key SKB, with out seeing the underlying plaintext, SKA, and SKB. let G be a multiplicative institution of high order q, and g be a random generator of G.

### 3.2 Re-encryption Key Generation

To outsource the re-encryption key technology system, we divide the re-encryption key era process into three one-of-a-kind levels.

First phase should be done by an originator  $U_i$ . The  $U_i$  generates the re-encryption key for a user  $U_j$  as meaning of that the  $U_i$  allows the  $U_j$  to share the data. For the first phase, the  $U_i$  performs  $RkGen_{init}(sk)$ : It picks  $s \in \mathbb{R} \mathbb{Z}_q$  and computes  $g^{-s \cdot ski}$ . It outputs  $DK = g^{-s \cdot ski}$ .

In the second phase, the  $U_i$  requests partial re-encryption computation for a set of target users  $\{U_j\}_{0 \leq j \leq \ell}$ , where  $\ell$  is number of users in sharing group, with sending the DK. Then outsourcing server performs  $RkGen_{partj}(pk_j, DK) = \{pkH3(g^{-s \cdot ski})^j, DKH3(pk_j)\}$  to generate a set of partial re-encryption keys,  $\{RK_{partj}\}_{0 \leq j \leq \ell}$ . In this point, the outsourcing server computes some partial computation that needs heavy exponential operation.

Third, after receiving  $\{RK_{partj}\}_{0 \leq j \leq \ell}$ , The  $U_i$  performs  $RkGen(RK_{partj})$  to complete encryption keys.

$$RK_{(i \rightarrow j)} = (rk1, rk2)$$

$$rk1 = (pk3^{(ski)^j})_s \quad rk2 = g^{-s \cdot ski \cdot 3(pk_j)}$$

Which makes dating between sub re-encryption keys for the robustness against selected encrypted attack. these schemes are based on n-Quotient Bilinear Diffie-Hellman Assumption, and the schemes proved protection by using calculating complexity of DBDH. inspired from the schemes, we designed the re-encryption key inclusive of values (rk1, rk2). because the rk1 consists of  $H3(rk2)$  and rk2 includes  $H3(pk_j)$ , it makes dating between rk1 and rk2. in this manner, the proposed scheme has robustness in opposition to chosen encrypted attack.

The start, in standard CPRE schemes, an originator desires to get hold of a hard and fast of public keys of users to make re-encryption key. The originator in our scheme wishes to pre-computed partial re-encryption key in place of the general public keys. consequently, most effective sending the DK is calculated as communication overhead.

### 3.3 Data Encryption

A CPRE scheme has two types of encryption.

$Enc1(m, pk_i)$ : It picks  $R \in \mathbb{R} \mathbb{G}$ . It computes  $r = H1(m, R)$ , and generates the first level encrypted  $CT = (C1, C2, C3, C4)$ ,  $C1 = g^r$ ,  $C2 = R \cdot e(g, pk_i)^{s \cdot H3(pk_i^{-s})}$ ,  $C3 = m \oplus H4(R)$ ,  $C4 = g^{s \cdot H3(-s \cdot ski) \cdot H3(pk_i)}$

$Enc2(m, w, pk_i)$ : It picks  $R \in \mathbb{R} \mathbb{G}$ . It computes  $r = H1(m, R)$ , and generates the second level encrypted  $CT_i = (C1, C2, C3, C4)$ ,  $C1 = g^r$ ,  $C2 = R \cdot e(pk_i, gs)W$ ,  $C3 = m \oplus H4(R)$ ,  $C4 = g^s \cdot H3(-s \cdot ski)$ ,

where  $W = H1(w || H4(C4))$ . 5.5.

### 3.4 Data Re-encryption

The outsourcing server performs  $ReEnc(CT_i, RK_{i \rightarrow j}, w)$  to re-encrypt the  $CT_i$  by  $U_j$ 's request. It computes  $CT_j$  as follows:

$$CT_j = (C1, C2, C3, C4)$$

$$C1 = C1,$$

$$C2 = R \cdot C2 \cdot e(g, rkW^2 \cdot rk1)$$

$$= R \cdot e(g, g)^{s \cdot ski \cdot H3(pk_j)} \cdot W \cdot e(g, g^{-s \cdot ski} \cdot W \cdot pks \cdot H3(g^{-s \cdot ski})^j)$$

$$= R \cdot e(g, pks \cdot H3(g^{-s \cdot ski}) \cdot H3(pk_j))^j$$

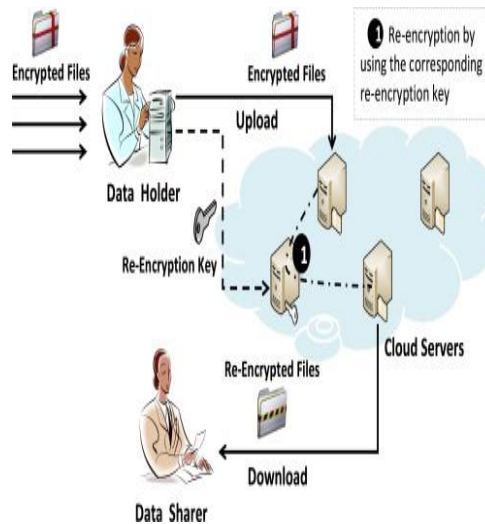
$$C3 = C3 = m \oplus H4(R),$$

$$C4 = C4 = g^s \cdot H3(-s \cdot ski)$$

Through the re-encryption process, a second level encrypted will be transformed to a first level encrypted of the user  $U_j$ .

**3.5 Encrypt:**

$N_{Z^+}$  is the number of conjunctive clauses in  $A$ ,  $n_i$  is the number of attributes in the  $i$ -th conjunctive clause  $CC_i$ , and  $a_{ij}$  is the  $j$ -th attribute in  $CC_i$ . Let  $DM_{i_1, \dots, i_t}$  be the DM administering all attributes in  $CC_i$ , where  $ID_{i_k}$  for  $1 < k < t$  are IDs of  $DM_{i_1, \dots, i_t}$ 's ancestors, and  $ID_{i_1} \neq ID_1$  is ID of the DM at the first level in a domain.



**Fig1.2:RE ENCRYPTED MODEL**

**IV. Analysis and Report**

In this work based on this set of rules like characteristic based totally proxy ReEncryption(ABPRE) used for this algorithm specially targeted on degree via stage. So we used information encrypted to accumulate the encrypted to accumulate the facts sequentially information encrypted

**V. CONCLUSION**

Attributed based totally proxy re-encryption scheme is appropriate on the cloud environment, because ABPRE can allow the information proprietor delegate the re-encryption right to the cloud for data sharing, and the records proprietor doesn't constantly be on-line. we proposed a HABE model, a HABE scheme, and a revocation mechanism, so that it will simultaneously attain: (1) excessive performance; (2) excellent-grained get admission to manage; (3) scalability; and (four) full delegation, in cloud computing. We describe the scalable revocation scheme from the systematic factor of view, and prove the HABE scheme, which is likewise collusion resistant, to be semantically cozy towards adaptive selected plaintext assaults beneath the BDH assumption and the random oracle version.

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